

Narrowing the Gap between TEEs Threat Model and Deployment Strategies

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Motivation Setting



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Sophisticated hardware layer attacks are possible

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Motivation Setting



Problematic, especially in malicious setting of blockchain (Maximal Extractable Value, ...)

- Compromising a **TEE**, could lead to large financial losses
- Not limited to blockchains AI model&data, ...

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Motivation Problem Definition

Trust assumptions are now on the host/operator being honest, to ensure physical security

However, current **TEE attestation flows** do not provide guarantees they operate in a given infrastructure

→ Provide assurance that CVM runs in the respective (trusted) infrastructure

How do we extend the attestation flow so the CVM runs on the respective infrastructure?

Background Deployments

Two deployments:

1. Bare/CVM flow





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Background Deployments

Two deployments:

- 1. Bare/CVM flow
- 2. Paravirtualization flow
 - Virtual Trust Level (VTL) for Intel or Virtual Machine Protection Level (VMPL) for AMD
 - Direction towards open source OpenHCL and COCONUT





Background Intel TDX

Intel TDX relies on two Intel SGX enclaves as a part of its attestation flow

- PCE (Provisioning Certificate Enclave)
- TDQE (Trust Domain Quoting Enclave)

In collaboration with Intel, receive an attestation key

As a part of the flow, provide PPID \rightarrow unique identifier of the platform



Design Use PPID?

We can rely on **PPID**

- Is part of the attestation flow
- Unique per platform
- Fixed for the platform

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What about the **binding** to the infrastructure provider?

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What about the **binding** to the infrastructure provider?

- **Provider** can create a database storing the values
- Verifier can query the database and receive True/False as an output

Design Use PPID – Challenges

Introduces some limitations:

- **Relies** on provider to share those information
 - Might vary across **cloud** providers
 - Extension to many parties, otherwise hard to integrate
- Limited visibility for the case of different CVMs on the same node
- Other **TEE** implementations
- Bare metal deployments?

Extension to Bare Metal Future directions

Two scenarios:

- 1. Confidential Virtual Machine (CVM) in cloud
- 2. Bare metal in cloud



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Summary Overview

Identification of the gap between attestation and threat model

Suggestion to strengthen it using **PPID** (or similar for AMD)

Poses several challenges

Future work should expand on the bare metal and less involvement of the provider

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Thank you!

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